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1: *****
2: maXbox Starter 47
3: *****
4:
5:
6: Work with real big RSA Cryptography
7: -----
8: Max Kleiner
9:
10: //Zwei Worte werden Dir im Leben viele Türen öffnen - "ziehen" und "stossen".
11:
12: A student of mine wrote a while ago asking me to clarify exactly how people used
    public key sharing to send encrypted messages. I had read a code a lot and thought
    I understood the concepts, but I cant explain it to with out a refresher.
13:
14: In its simplest forms, an individual has public and private keys, each including
15: of two values.
16: For the public key the values are n of p*q, the so called "modulus", and E, a well
    known encrypting integer prime with the value: Const E = 65537;.
17:
18: The private key values are also n, the same modulus that appears in the public key,
    and d, a big number which can decrypt any message encrypted using the public key.
19:
20: There are obviously two cases:
21:
22:     1. Encrypting with the public key, and then decrypting with the private key.
23:         For a message or data
24:     2. Encrypting with the private key, and then decrypting with the public key.
25:         For a digital signature
26:
27: Lets begin with a big real world example:
28:
29:     1. First we have to generate two very large primes, a(p) and a(q). These numbers
    must be random and not too close to each other. Here are the numbers that we
    generated by using Rabin-Miller primality tests:
30:
31: Const
32:     ap =
33: '1213107243921127189732367153161244042847242763370141092563454931230196437304208
34: '5619324197365322416866541017057361365214171711713797974299334871062829803541';
35:     aq =
36: '1202752425547874888595622079373451212873338780368207543365389998395517985098879
37: '7899869146900809131611153346817050832096022160146366346391812470987105415233';
38:
39: The number is of a RSA crypto system with 1024 bits (2^1024), so the big number n
    (modulus)
40: has a length of 309:
41:
42: maxcalcF('ln(2^512)/ln(10)') >>> 154.127357779958 //one prime
43: maxcalcF('ln(2^1024)/ln(10)') >>> 308.254715559917 //modulus
44:
45:     2. Second with these two large prime numbers, we can calculate n and phi(n):
46:
47: n = ap * aq
48: phi(n) = (ap - 1) * (aq - 1)
49:
50: The RSA generation technique depends on the fact that it is generally difficult to
    factor a number which is the product of two "large" prime numbers. This is the way
    that RSA determines the n value!
51:
52: n is >>> writeln(bigMulu(ap,aq)) =
53:
54:     145906768007583323230186939349070635292401872375357164399581871019873438799005
55:     358938369571402670149802121818086292467422828157022922076746906543401224889672
56:     472407926969987100581290103199317858753663710862357656510507883714297115637342
57:     788911463535102712032765166518411726859837988672111837205085526346618740053
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59: phi(n) is >>> writeln(bigMulu(ap-1,aq-1)) =
60:
61:     145906768007583323230186939349070635292401872375357164399581871019873438799005
62:     358938369571402670149802121818086292467422828157022922076746906543401224889648
63:     313811232279966317301397777852365301547848273478871297222058587457152891606459
64:     269718119268971163555070802643999529549644116811947516513938184296683521280
65:
66: The Euler phi-function assigns to an integer n the number of numbers less than n
and relatively prime to n. For example phi(12)= 4. Why?
67: The only numbers less than 12 which are relatively prime to 12 are 1,5,7,11 so
phi(12) = 4.
68:
69:     writeln('EulerPhi '+itoa(GetEulerPhi(12)));
70:     >>> 4
71:
72: What about a prime as phi(11)?
73: Since 11 is prime, all the numbers less than 11 are relatively prime to 11, so
phi(11) = 10.
74:
75:     writeln('EulerPhi '+itoa(GetEulerPhi(11)));
76:     >>> 10
77:
78: This can be checked by counting the number of numbers less than N which have a
factor in common with N by Eulers Theorem of phi of n.
79:
80: 3. Third we have to find the private Key, the difficult part of RSA.
81: Euclid gives a process which we now call the Euclidean Algorithm. This process,
essentially repeated long division, can be used (The Extended Euclidean Algorithm)
to find the private key:
82:
83: //Find d such that: e*d mod phi(n) = 1. (as inverse modulo)
84: function getPrivate: integer;
85: begin
86:   for it:= 1 to big(1200) do //it(n) depends on bigint
87:     if (getEPublic * it mod getPHI_N = 1) then begin
88:       result:= it;
89:       privKey:= result;
90:       break;
91:     end
92:   end;
93:
94: This multiplicative inverse is the private key [e*d=1 mod phi(n)]. The common
notation for expressing the private key is d or in our case addpkey:
95:
96: Const addpkey = // - the private key -
97: 89489425009274444368228545921773093919669586065884257445497854456487674839629818
98: 39093494197326287961679797060891728367987549933157416111385408881327548811058824
99: 71930775825272784379065040156806234235500672400424666656542323835029222154936232
100: 89472138866445818789127946123407807725702626644091036502372545139713';
101:
102: Now we do the proof of the crypto pudding:
103:
104: Find d such that: e*d mod phi(n)=1
105:
106:   Extended Euclidean Algorithm with real world example e*d mod phi(n)=1
107:
108:   writeln(modbig2(bigmulu(addpkey,65537),bigmulu(ap1,aq1)));
109:   >>> 1
110:
111:   Or with a smaller example:
112:   e*d mod phi(n)=1 >>> 7*103=721 mod 120 =1. [p=11 and q=13.] //invmod()
113:   maXcalcF('7 * 103 % 120') >>> 1
114:
115: 4. Fourth we are able to cipher data with:
116:
117:   encrypt: m^e mod n = 9^7 mod 143 = 48 = c; //9 is the message
118:   decrypt: c^d mod n = 48^103 mod 143 = 9 = m;

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119:     below: Crypto Systems Ref 2
120:
121: Our big encrypt: m^e mod n, m= 1976620216402300889624482718775150
122:
123: writeln(modpowbig2('1976620216402300889624482718775150',65537,bigmulu(ap,aq)))
124: s_encrypt:= modpowbig2('1976620216402300889624482718775150',65537,bigmulu(ap,aq));
125: s_decrypt:= modpowbig2(s_encrypt,addpkey,bigmulu(ap,aq))
126: writeln('decrypt: '+s_decrypt)
127:
128: RSA is the single most useful applied concept for building cryptographic
129: protocols.
130:
131: Ref: http://www.softwareschule.ch/maxbox.htm
132:
133: ..\examples\210_RSA_crypto_complete8hybrid.txt
134: ..\examples\750_ibz_cryptomem_RSA_proof_64.txt
135:
136:     Crypto Systems Ref:
137: 1. Prime1: 13 and Prime2: 23
138: 2. RSA Modul Public [p*q] as N: 299
139: 3. phi(N) Private: 264 = 12*22
140: 4. RSA Exponent: 17
141: 5. Private D: 233
142: 6. Public (17,299) - Private (233,299) //}
143:
144:     Crypto Systems Ref 2:
145: 1. Prime1: 11 & Prime2: 13
146: 2. RSA Modul Public [p*q] as N: 143
147: 3. Phi(N) Private: 120 = 10*12
148: 4. Public RSA Exponent: 7
149: 5. Private D: 103
150: 6. Public (7,143) - Private (103,143)
151:
152: Lets choose two primes: p=11 and q=13.
153: Hence the modulus is n=p*q=143. The totient of n phi(n)=(p-1)*(q-1)=120.
154:
155: Doc:
156: http://doctrina.org/How-RSA-Works-With-Examples.html
157: http://www.delphiforfun.org/Programs/Math_Topics/RSA_KeyDemo.htm
158: http://www.softwareschule.ch/download/maxbox_functions.txt
159:
160: There are only 10 types of people: those who understand binary and those who
161: don't.
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